

Data Communications and Networking Fourth Edition



Multiple Access

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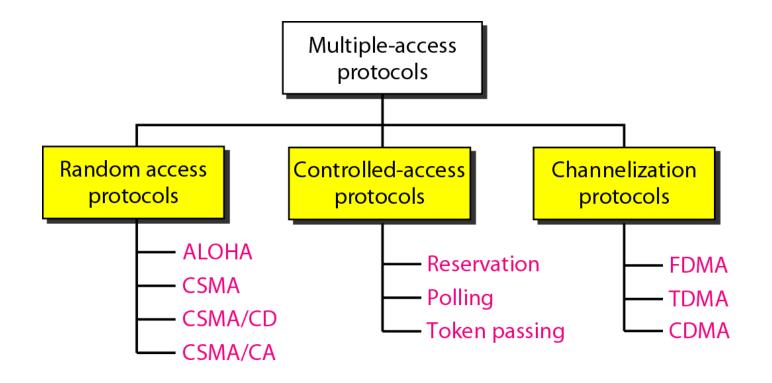
Data link layer divided into two functionality-oriented sublayers

Data link layer

Data link control

Multiple-access resolution

Taxonomy of multiple-access protocols discussed in this chapter



In random access or contention methods, no station is superior to another station and none is assigned the control over another. No station permits, or does not permit, another station to send. At each instance, a station that has data to send uses a procedure defined by the protocol to make a decision on whether or not to send.

Topics discussed in this section: ALOHA Carrier Sense Multiple Access Carrier Sense Multiple Access with Collision Detection Carrier Sense Multiple Access with Collision Avoidance To avoid access conflict or to resolve it when it happens, each station follows a procedure that answer the following questions:

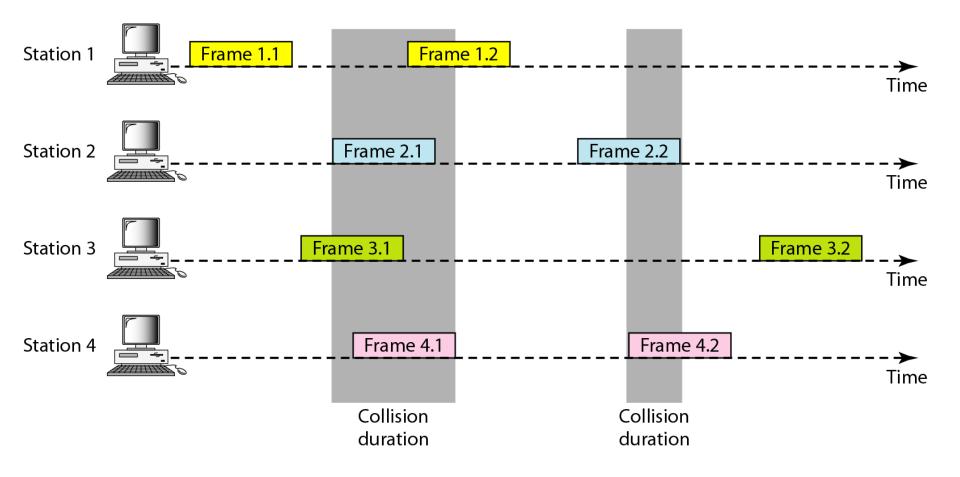
When can the station access the medium.
What can the station do if the medium is busy.
How can the station determine the success or failure of the transmission.

What can the station do if there is an access conflict.

ALOHA, the earliest random access method, was developed at the University of Hawaii in early 1970. It was designed for a radio (wireless) LAN, but it can be used on any shared medium.

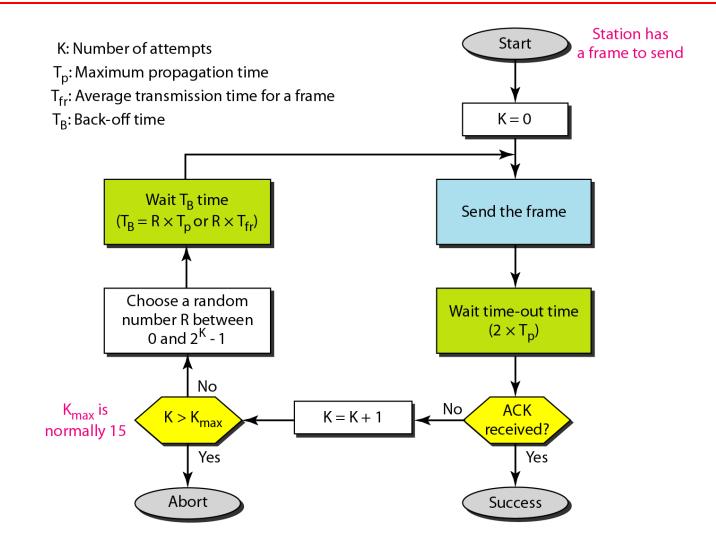
Pure ALOHA Slotted ALOHA

Frames in a pure ALOHA network



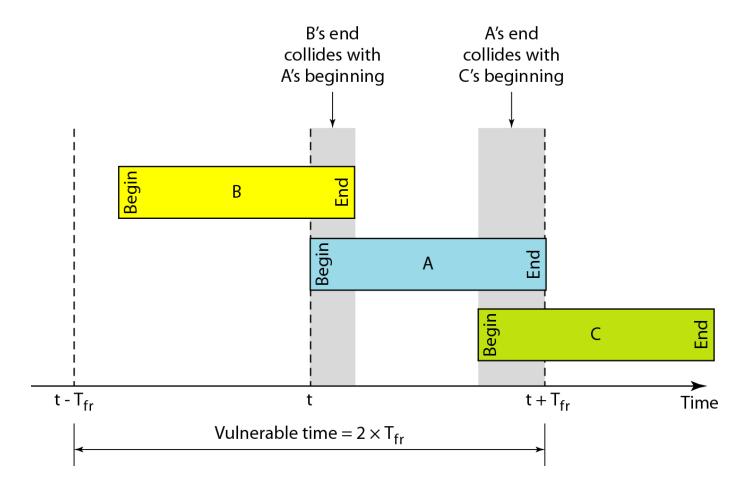
12.7

Procedure for pure ALOHA protocol



Vulnerable time for pure ALOHA protocol

The Vulnerable time, in which there is possibility of collision



Slotted ALOHA

Frames in a slotted ALOHA network

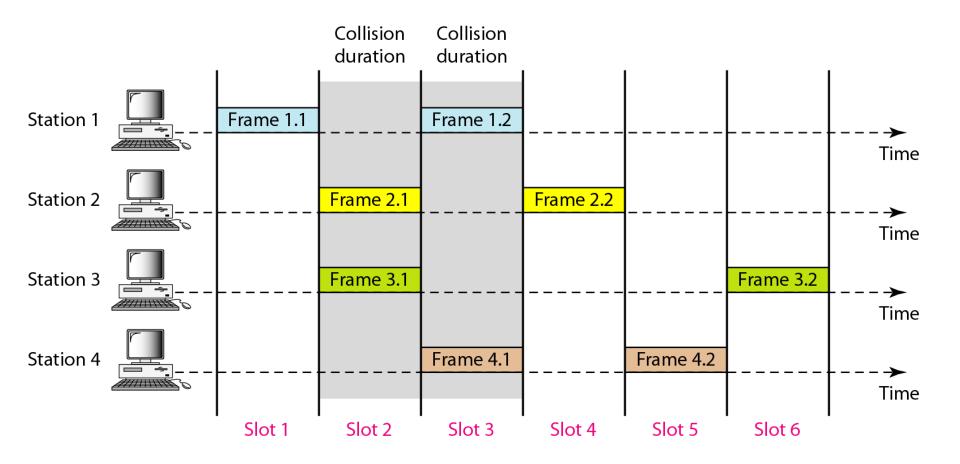
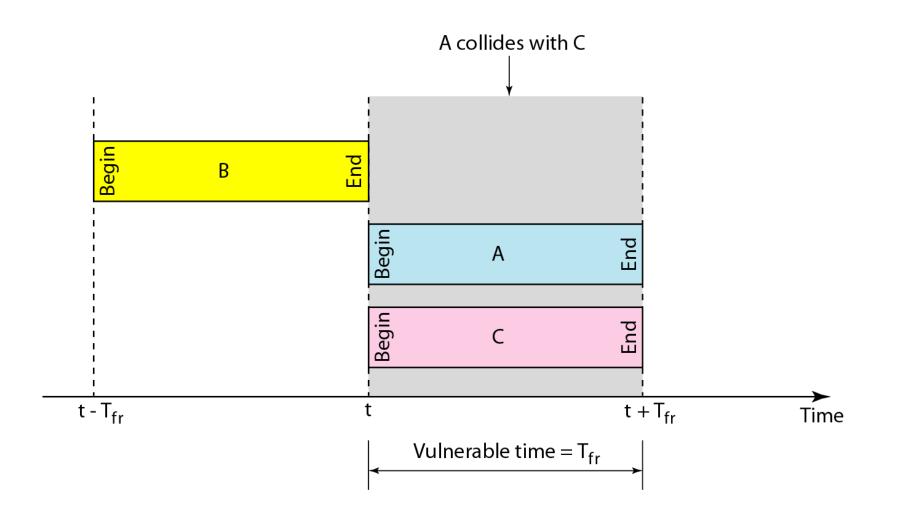


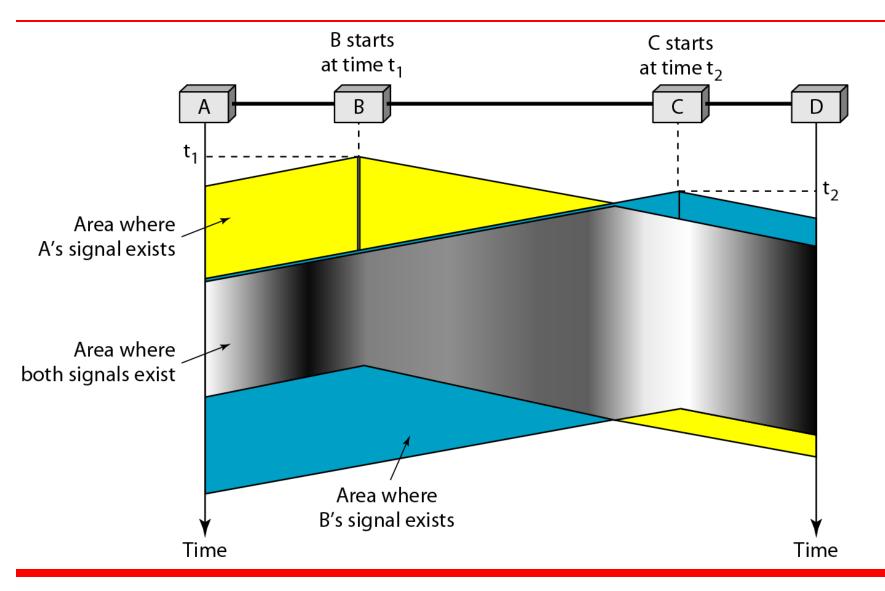
Figure 12.7 Vulnerable time for slotted ALOHA protocol



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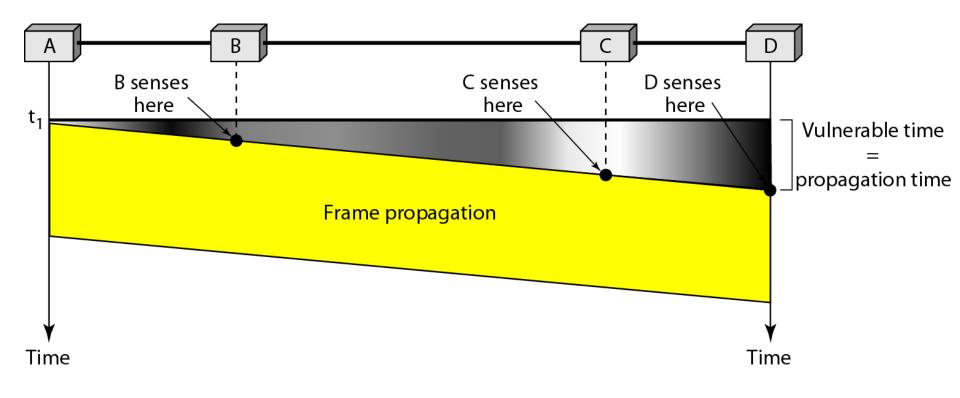
CSMA Carrier Sense Multiple Access

Space/time model of the collision in CSMA



12.14

Figure 12.9 *Vulnerable time in CSMA*

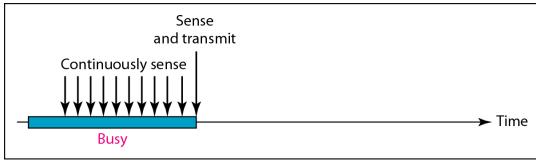


Persistence Method

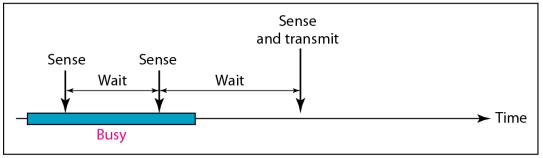
What Should a station do if the channel is busy. What should a station do if the channel is idle.

1-Persistence Method Non Persistence Method p-Persistence Method

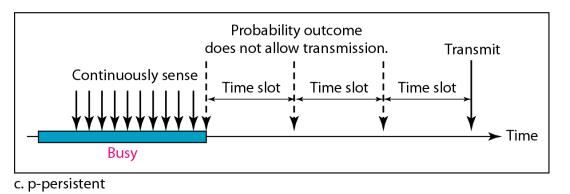
Behavior of three persistence methods



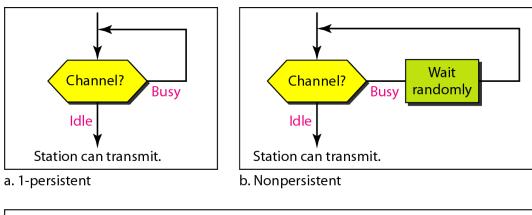
a. 1-persistent

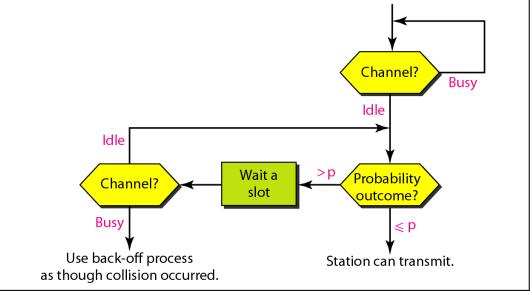


b. Nonpersistent



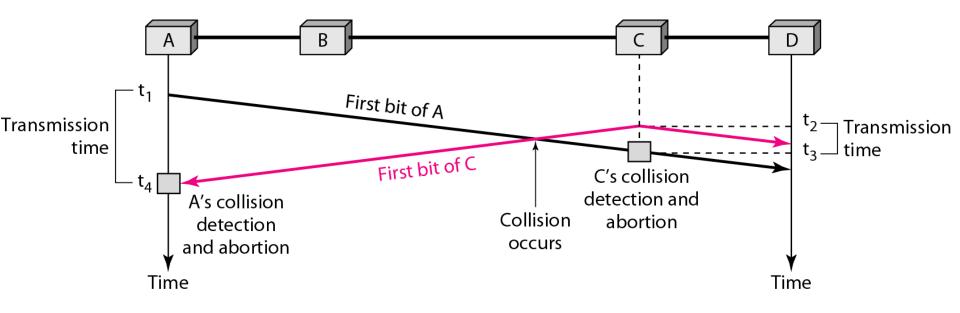
Flow diagram for three persistence methods



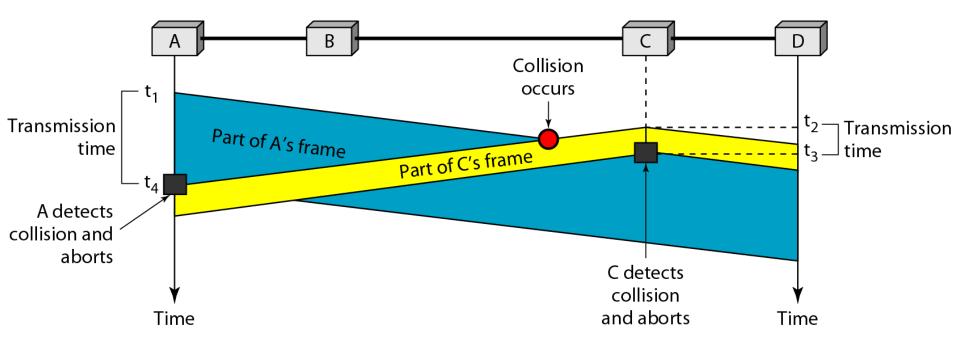


c. p-persistent

Collision of the first bit in CSMA/CD



Collision and abortion in CSMA/CD



Flow diagram for the CSMA/CD

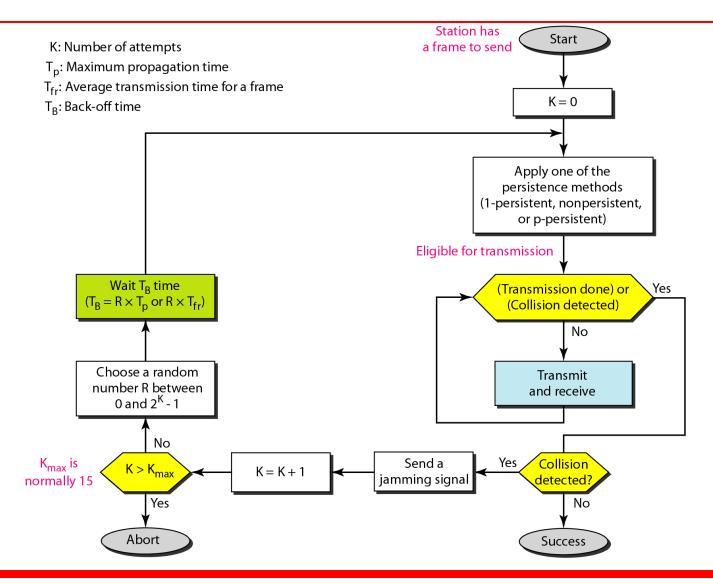


Figure 12.15 Energy level during transmission, idleness, or collision

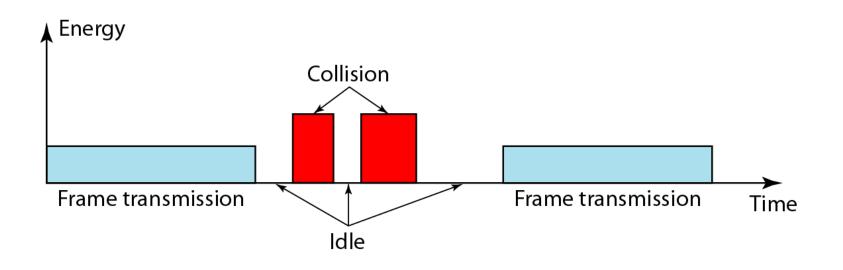
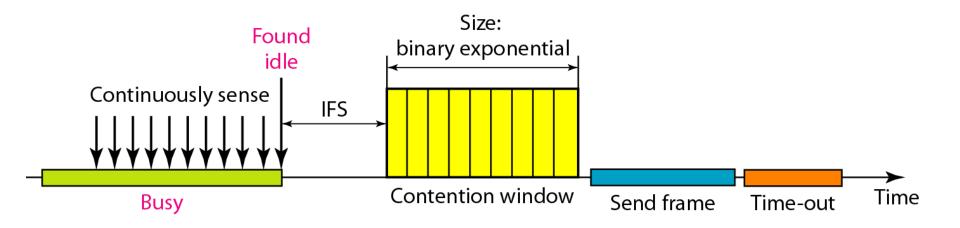


Figure 12.16 *Timing in CSMA/CA*

Interframe Space (IFS)



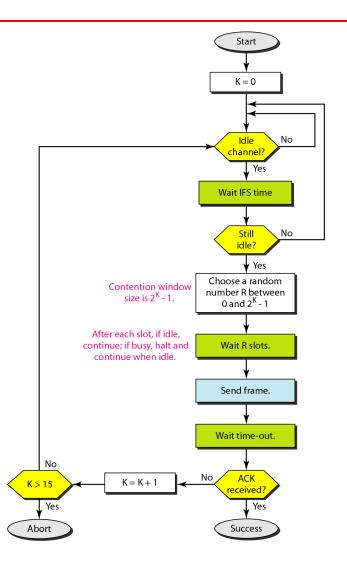


In CSMA/CA, the IFS can also be used to define the priority of a station or a frame.



In CSMA/CA, if the station finds the channel busy, it does not restart the timer of the contention window; it stops the timer and restarts it when the channel becomes idle.

Figure 12.17 Flow diagram for CSMA/CA



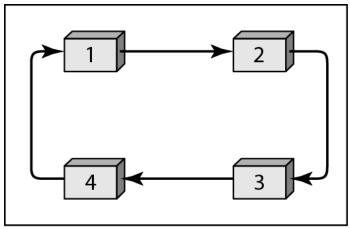
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CONTROLLED ACCESS

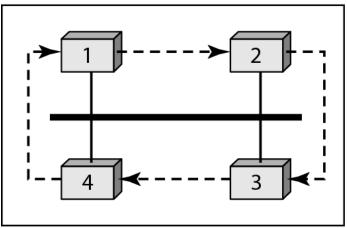
In controlled access, the stations consult one another to find which station has the right to send. A station cannot send unless it has been authorized by other stations. We discuss three popular controlled-access methods.

Topics discussed in this section: Reservation Polling Token Passing

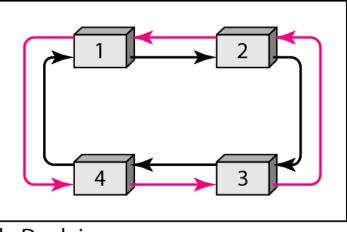
Logical ring and physical topology in token-passing access method



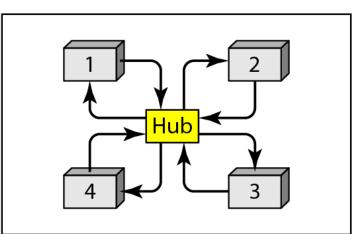
a. Physical ring



c. Bus ring



b. Dual ring



d. Star ring

DRONACHARYA COLLEGE OF ENGINEERING, GURGAON Computer Science & Engineering Assignment Semester- VI (I & II) Subject with Code: Computer Networks (IT-305-F)

Q:1 Explain Lacal Area Network. Q:2 Explain IEEE 802 standards in Detail.